Pharo VM Runtime























VM C Runtime

Compiled Ahead of Time Slang -> C -> Native code





VM C Runtime

Compiled Ahead of Time Slang -> C -> Native code



VM Generated Runtime

Compiled at Run Time Bytecode -> Native code



















GIGUENATE RUNTIME PUN Confiled Bend if possible Send Stay in native Send Code mon-Capilad

How they interact at run time Private Statige Bridge GIGUENATED RUNTIME interpret comple PUN Confiled method Bend if possible Send Stay in native Code Fun mon-Capilad Dend interpretent possible







VM C Runtime vs Generated Runtime

- The JITted methods do not include code for all cases
- Slow paths, complex cases (such as the GC) are NOT compiled

Instead => delegate to the VM C code ... But...









Trampolines vs Reverse Trampolines Two kinds of trampolines

- Trampolines:
 - Generated native code —> C code

- Reverse trampolines:
 - C code —> Generated native code

What are trampolines used for? Mostly, slow execution paths

- Send a message to a non-JITted method
- PICs
- Call the GC
- Slow allocations
- Immutability checks
- Managing GC invariants (e.g., remembered set)
- Must be boolean

What are reverse trampolines used for?

- Call JITted methods
- Low level Routines (e.g., get the state of the machine)

That's most of it

Why do we need trampolines? Communication Problems

- The C functions where compiled by **some** C compiler
 - We don't know: Are they optimised? What registers do they use?
 - We know: they have a standard calling convention
 - We have to be conservative and consider C functions could do anything!
- Our generated native code
 - Does not follow the C calling convention
 - Runs in a separate stack!

Trampolines Calling the VM C runtime from the Generated Runtime

- Native code routines that make the transition
- Jitted methods call a trampoline
- The trampoline sets-up the conditions to call C code
- Then calls the C code



compoline? 07 call C function!



Two Execution Stacks One for the Pharo execution, one for the VM C runtime execution

- The Pharo execution is persistent (e.g., when you save the image)
- And has green threads!
- The C execution is single-threaded non-persistent

=> Two stacks to **separate** concerns



Cstuck main(interport

Trampolines in the Pharo VM Switching to the C stack



are trapplines prepare wys switch to C stack call C function!



Trampolines in the Pharo VM Switching to the C stack



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Reverse Trampolines (enilopmarts) Calling the generated runtime from C runtime

- Native code routines that make the transition
- The VM C runtime calls a reverse trampoline
- The reverse trampoline sets-up the conditions to call the generated runtime
- Then calls the generated code







Reverse Trampolines in the Pharo VM Switching to the Pharo Stack





Reverse Trampolines in the Pharo VM Switching to the Pharo Stack





Where are the trampolines? The Native Code Zone

- It has Intrinsics + Jitted Methods
- Intrinsics: routines that are predefined by the compiler
 - Intrinsics are generated at VM startup and stay there
 - Trampolines are not the only intrinsics (e.g., marrying a machine code frame to a context object)



Other details about trampolines

- Reverse trampolines do not actually call, they return to generated code, to avoid having the reverse trampoline in the stack
- Reverse trampolines do not receive arguments, arguments are passed through the smalltalk stack
 - They then pop the arguments and puts them in the right place before returning to the generated method
- Pharo Trampolines do not ensure that registers are properly saved-restored. It's the caller method that ensures that. This is specially problematic since the called C function could do virtually **anything**, even smashing our registers!



